

## Relational Algebra

Kuang-hua Chen  
Department of Library and Information Science  
National Taiwan University

## Operations on Relational Model

- Definition
- Updates
  - Insert
  - Delete
  - Modify
- Retrievals
  - Relational algebra

Language & Information Processing System, LIS, NTU

2

## Define Relational Databases

```

DECLARE SCHEMA COMPANY;
DECLARE DOMAIN PERSON_SSNS TYPE FIXED_CHAR (9);
DECLARE DOMAIN PERSON_NAMES TYPE
    VARIABLE_CHAR(15);
DECLARE DOMAIN PERSON_INITIALS TYPE
    ALPHABETIC_CHAR(1);
DECLARE DOMAIN DATES TYPE DATE;
DECLARE DOMAIN ADDRESSES TYPE VARIABLE_CHAR(35);
DECLARE DOMAIN PERSON_SEX TYPE ENUMERATED {M, F};
DECLARE DOMAIN PERSON_SALARIES TYPE MONEY;
DECLARE DOMAIN DEPT_NUMBERS TYPE INTEGER_RANGE
    [1,10];
DECLARE DOMAIN DEPT_NAMES TYPE VARIABLE_CHAR(20);
    
```

Language & Information Processing System, LIS, NTU

3

## Define a Relation

```

DECLARE RELATION EMPLOYEE
FOR SCHEMA COMPANY
ATTRIBUTES FNAME          DOMAIN PERSON_NAMES,
             MINIT         DOMAIN PERSON_INITIALS,
             LNAME         DOMAIN PERSON_NAMES,
             SSN           DOMAIN PERSON_SSNS,
             BDATE        DOMAIN DATES,
             ADDRESS      DOMAIN ADDRESSES,
             SEX          DOMAIN PERSON_SEX,
             SALARY       DOMAIN PERSON_SALARIES,
             SUPERSSN     DOMAIN PERSON_SSNS,
             DNO          DOMAIN DEPT_NUMBERS

CONSTRAINTS PRIMARY_KEY (SSN),
             FOREIGN_KEY (SUPERSSN) REFERENCES EMPLOYEE,
             FOREIGN_KEY (DNO) REFERENCES DEPARTMENT;
    
```

Language & Information Processing System, LIS, NTU

4

## Define another Relation

```

DECLARE RELATION DEPARTMENT
FOR SCHEMA COMPANY
ATTRIBUTES DNAME          DOMAIN DEPT_NAMES,
             DNUMBER       DOMAIN DEPT_NUMBERS,
             MGRSSN        DOMAIN PERSON_SSNS,
             MGRSTARTDATE DOMAIN DATES

CONSTRAINTS PRIMARY_KEY (DNUMBER),
             KEY (DNAME),
             FOREIGN_KEY (MGRSSN) REFERENCES EMPLOYEE;
    
```

Language & Information Processing System, LIS, NTU

5

## The COMPANY Relational Database Schema

```

EMPLOYEE
[FNAME | MINIT | LNAME | SSN | BDATE | ADDRESS | SEX | SALARY | SUPERSSN | DNO]

DEPARTMENT
[DNAME | DNUMBER | MGRSSN | MGRSTARTDATE]

DEPT_LOCATIONS
[DNUMBER | DLOCATION]

WORKS_ON
[ESSN | PNO | HOURS]

PROJECT
[PNAME | PNUMBER | PLOCATION | DNUM]

DEPENDENT
[ESSN | DEPENDENT_NAME | SEX | BDATE | RELATIONSHIP]
    
```

Language & Information Processing System, LIS, NTU

6

## A Relational Database Instance of the COMPANY Schema

EMPLOYEE	FNAME	MINIT	LNAME	SSN	BDATE	ADDRESS	SEX	SALARY	SUPERSSN	DNO
John	B	Smith		123456789	09-JAN-55	731 Fondren, Houston, TX	M	30000	333445555	5
Franklin	T	Wong		333445555	08-DEC-45	638 Voss, Houston, TX	M	40000	888665555	5
Alicia	J	Zelaya		999887777	19-JUL-58	3321 Castle, Spring, TX	F	25000	987654321	4
Jennifer	S	Wallace		987654321	20-JUN-31	291 Berry, Bellaire, TX	F	43000	888665555	4
Ramesh	K	Narayan		666884444	15-SEP-52	975 Fire Oak, Humble, TX	M	38000	333445555	5
Joyce	A	English		453453453	31-JUL-62	5631 Rice, Houston, TX	F	25000	333445555	5
Ahmad	V	Jabbar		987987987	29-MAR-59	980 Dallas, Houston TX	M	25000	987654321	4
James	E	Borg		888665555	10-NOV-27	450 Stone, Houston, TX	M	55000	null	1

## A Relational Database Instance of the COMPANY Schema (Continued)

DEPT_LOCATIONS	DNUMBER	DLOCATION
	1	Houston
	4	Stafford
	5	Bellaire
	5	Sugarland
	5	Houston

DEPARTMENT	DNAME	DNUMBER	MGRSSN	MGRSTARTDATE
Research		5	333445555	22-MAY-78
Administration		4	987654321	01-JAN-85
Headquarters		1	888665555	19-JUN-71

## A Relational Database Instance of the COMPANY Schema (Continued)

WORKS_ON	ESSN	PNO	HOURS	PROJECT	PNAME	PNUMBER	PLOCATION	DNUM
	123456789	1	32.5	ProductX	1	Bellaire	5	
	123456789	2	7.5	ProductY	2	Sugarland	5	
	666884444	3	40.0	ProductZ	3	Houston	5	
	453453453	1	20.0	Computerization	10	Stafford	4	
	453453453	2	20.0	Reorganization	20	Houston	1	
	333445555	2	10.0	Newbenefits	30	Stafford	4	
	333445555	3	10.0					
	333445555	10	10.0					
	333445555	20	10.0					
	999887777	30	30.0					
	999887777	10	10.0					
	987987987	10	35.0					
	987987987	30	5.0					
	987654321	30	20.0					
	987654321	20	15.0					
	888665555	20	null					

## A Relational Database Instance of the COMPANY Schema (Continued)

DEPENDENT	ESSN	DEPENDENT_NAME	SEX	BDATE	RELATIONSHIP
	333445555	Alicia	F	05-APR-076	DAUGHTER
	333445555	Theodore	M	25-OCT-73	SON
	333445555	Joy	F	03-MAY-48	SPOUSE
	987654321	Abner	M	29-FEB-32	SPOUSE
	123456789	Michael	M	01-JAN-78	SON
	123456789	Alice	F	31-DEC-78	DAUGHTER
	123456789	Elizabeth	F	05-MAY-57	SPOUSE

## Insert

- Insert <'Cecilia', 'F', 'Kolonsky', '677678989', '05-Apr-50', '6357 Windy Lane, Katy, TX', F, 28000, null, 4> into EMPLOYEE
- Insert <'Alicia', 'J', 'Zelaya', '999887777', '05-Apr-50', '6357 Windy Lane, Katy, TX', F, 28000, '987654321', 4> into EMPLOYEE
  - violate the key constraint
- Insert <'Cecilia', 'F', 'Kolonsky', null, '05-Apr-50', '6357 Windy Lane, Katy, TX', F, 28000, null, 4> into EMPLOYEE
  - violate the entity integrity

## Delete

- Can violate only referential integrity
- Delete the WORKS\_ON tuple with PNO=10 and ESSN='999887777'
- Delete the EMPLOYEE tuple with SSN='999887777'
  - violate the referential integrity
  - reject the delete
  - cascade the deletion
  - modify the referencing attribute

## Modify

- Modifying an attribute which is neither a primary key nor foreign key causes no problems
- Modifying a primary key value is similar to deleting one tuple and inserting another
- Modifying a foreign key may violate referential integrity

## Operations of Relation Algebra

- Set operations
  - Union
  - Intersection
  - Difference
  - Product
- Born original operations
  - Select
  - Project
  - Join
  - Aggregate function

## RESULT ← RESULT1 ∪ RESULT2

RESULT1	SSN	RESULT	SSN
	123456789		123456789
	333445555		333445555
	666884444		666884444
	453453453		453453453
			888665555

  

RESULT2	SSN
	333445555
	888665555

## UNION, INTERSECTION, and DIFFERENCE

(a) Two union compatible relations.

STUDENT	FN	LN	INSTRUCTOR	FNAME	LNAME
	Susan	Yao		John	Smith
	Ramesh	Shah		Ricardo	Browne
	Johnny	Kohler		Susan	Yao
	Barbara	Jones		Francis	Johnson
	Amy	Ford		Ramesh	Shah
	Jimmy	Wang			
	Ernest	Gilbert			

## UNION, INTERSECTION, and DIFFERENCE (Continued)

(b) STUDENT ∪ INSTRUCTOR (c) STUDENT ∩ INSTRUCTOR

FN	LN	FN	LN
Susan	Yao	Susan	Yao
Ramesh	Shah	Ramesh	Shah
Johnny	Kohler		
Barbara	Jones		
Amy	Ford		
Jimmy	Wang		
Ernest	Gilbert		
John	Smith		
Ricardo	Browne		
Francis	Johnson		

## UNION, INTERSECTION, and DIFFERENCE (Continued)

(d) STUDENT - INSTRUCTOR (e) INSTRUCTOR - STUDENT

FN	LN	FNAME	LNAME
Jonny	Kohler	John	Smith
Barbara	Jones	Ricardo	Browne
Amy	Ford	Francis	Johnson
Jimmy	Wang		
Ernest	Gilbert		

## CARTESIAN PRODUCT

FEMALE_EMPS	FNAME	MINIT	LNAME	SSN	BDATE	ADDRESS	SEX	SALARY	SUPERSSN	DWO
Alicia	J	Zelaya	999887777	19-JUL-58	3321 Castle, Spring, TX	F	25000	987654321	4	
Jennifer	S	Wallace	987654321	20-JUN-31	291 Berry, Bellaire, TX	F	43000	888665555	4	
Joyce	A	English	453453453	31-JUL-62	5631 Rice, Houston, TX	F	25000	333445555	5	

EMP_NAMES	FNAME	LNAME	SSN	RESULT	FNAME	LNAME	DEPENDENT_NAME
Alicia	Zelaya	999887777		Jennifer	Wallace	Abner	
Jennifer	Wallace	987654321					
Joyce	English	453453453					

ACTUAL_DEPENDENTS	FNAME	LNAME	SSN	ESSN	DEPENDENT_NAME	SEX	BDATE
	Jennifer	Wallace	987654321	987654321	Abner	M	29-FEB-32

## CARTESIAN PRODUCT (Continued)

EMP_DEPENDENTS	FNAME	LNAME	SSN	ESSN	DEPENDENT_NAME	SEX	BDATE	...
Alicia	Zelaya	999887777	333445555	Alice	F	05-APR-76	...	
Alicia	Zelaya	999887777	333445555	Theodore	M	25-OCT-73	...	
Alicia	Zelaya	999887777	333445555	Joy	F	03-MAY-48	...	
Alicia	Zelaya	999887777	987654321	Abner	M	29-FEB-32	...	
Alicia	Zelaya	999887777	123456789	Michael	M	01-JAN-78	...	
Alicia	Zelaya	999887777	123456789	Alice	F	31-DEC-78	...	
Alicia	Zelaya	999887777	123456789	Elizabeth	F	05-MAY-57	...	
Jennifer	Wallace	987654321	333445555	Alice	F	05-APR-76	...	
Jennifer	Wallace	987654321	333445555	Theodore	M	25-OCT-73	...	
Jennifer	Wallace	987654321	333445555	Joy	F	03-MAY-48	...	
Jennifer	Wallace	987654321	987654321	Abner	M	29-FEB-32	...	
Jennifer	Wallace	987654321	123456789	Michael	M	01-JAN-78	...	
Jennifer	Wallace	987654321	123456789	Alice	F	31-DEC-78	...	
Jennifer	Wallace	987654321	123456789	Elizabeth	F	05-MAY-57	...	
Joyce	English	453453453	333445555	Alice	F	05-APR-76	...	
Joyce	English	453453453	333445555	Theodore	M	25-OCT-73	...	
Joyce	English	453453453	333445555	Joy	F	03-MAY-48	...	
Joyce	English	453453453	987654321	Abner	M	29-FEB-32	...	
Joyce	English	453453453	123456789	Michael	M	01-JAN-78	...	
Joyce	English	453453453	123456789	Alice	F	31-DEC-78	...	
Joyce	English	453453453	123456789	Elizabeth	F	05-MAY-57	...	

## The DIVISION Operation

SSN_PNOS	ESSN	PNO
123456789	1	
123456789	2	
666884444	3	
453453453	1	
453453453	2	
453453453	3	
333445555	4	
333445555	10	
333445555	20	
999887777	30	
999887777	10	
987987987	10	
987987987	30	
987654321	30	
987654321	20	
888665555	20	

Dividing SSN\_PNOS by SMITH\_PNOS.

$SSNS(SSN) \leftarrow SSN\_PNOS \div SMITH\_PNOS$

SMITH_PNOS	PNO
	1
	2

SSNS	SSN
	123456789
	453453453

## The DIVISION Operation (Continued)

R	A	B
	a1	b1
	a2	b1
	a3	b1
	a4	b1
	a1	b2
	a3	b2
	a2	b3
	a3	b3
	a4	b3
	a1	b4
	a2	b4
	a3	b4

(b)  $T \leftarrow R \div S$

S	A
	a1
	a2
	a3

T	B
	b1
	b4

## The DIVISION Operation (Continued)

- $T = R \div S$
- Assume  $Y = \text{att}[R] - \text{att}[S]$
- $T_1 = \prod_Y R$
- $T_2 = \prod_Y (S \times T_1 - R)$
- $T = T_1 - T_2$

## SELECT

- use to select a subset of the tuples in a relation that satisfy a selection condition
- $\sigma_{\langle \text{selection condition} \rangle} \langle \text{Table} \rangle$
- $\langle \text{selection} \rangle$  is made up of a number of clauses
  - $\langle \text{attribute name} \rangle \langle \text{comparison op} \rangle \langle \text{constant value} \rangle$
  - $\langle \text{attribute name} \rangle \langle \text{comparison op} \rangle \langle \text{attribute name} \rangle$
- The fraction of tuples selected by a selection condition is referred to as the selectivity of the condition

## PROJECT

- PROJECT extracts certain columns from the table and discards the other columns
- PROJECT is used to “project” attributes we are interested in
- $\pi$  <attribute list> (<relation name>)
- The resulting relation has only the attributes specified in <attribute list> and in the same order as they appear in the list
- The degree equals to the number of attributes in <attribute list>
- Duplicate elimination

## SELECT and PROJECT Operations

(a)  $\sigma$  (DNO=4 AND SALARY>25000) OR (DNO=5 AND SALARY>30000) (EMPLOYEE)

FNAME	MINIT	LNAME	SSN	BDATE	ADDRESS	SEX	SALARY	SUPERSSN	DNO
Franklin	T	Wong	333445555	08-DEC-45	638 Voss, Houston, TX	M	40000	888665555	5
Jennifer	S	Wallace	987654321	20-JUN-01	291 Berry, Bellare, TX	F	43000	888665555	4
Ramesh	K	Narayan	666884444	15-SEP-52	975 FireOak, Humble, TX	M	38000	333445555	5

(b)  $\pi$  LNAME, FNAME, SALARY (EMPLOYEE)

LNAME	FNAME	SALARY
Smith	John	30000
Wong	Franklin	40000
Zelaya	Alicia	25000
Wallace	Jennifer	43000
Narayan	Ramesh	38000
English	Joyce	25000
Jabbar	Ahmad	25000
Borg	James	55000

(c)  $\pi$  SEX, SALARY (EMPLOYEE)

SEX	SALARY
M	30000
M	40000
F	25000
F	43000
M	38000
M	25000
M	55000

## Renaming of Attributes

- $\pi$  LNAME, FNAME, SALARY ( $\sigma_{DNO=5}$  (EMPLOYEE))
- DEP5\_EMPS  $\leftarrow \sigma_{DNO=5}$  (EMPLOYEE)
- RESULT  $\leftarrow \pi$  LNAME, FNAME, SALARY (DEP5\_EMPS)
- R(LASTNAME, FIRSTNAME, SALARY)  $\leftarrow \pi$  LNAME, FNAME, SALARY (DEP5\_EMPS)

## Results of Relational Algebra

(a)  $\pi$  LNAME, FNAME, SALARY ( $\sigma_{DNO=5}$  (EMPLOYEE))

FNAME	LNAME	SALARY
John	Smith	30000
Franklin	Wong	40000
Ramesh	Narayan	38000
Joyce	English	25000

## Results of Relational Algebra (Continued)

(b) The same expression using intermediate relations and renaming of attributes.

TEMP	FNAME	MINIT	LNAME	SSN	BDATE	ADDRESS	SEX	SALARY	SUPERSSN	DNO
	John	B	Smith	123456789	09-JAN-55	731 Fondren, Houston, TX	M	30000	333445555	5
	Franklin	T	Wong	333445555	08-DEC-45	638 Voss, Houston, TX	M	40000	888665555	5
	Ramesh	K	Narayan	666884444	15-SEP-52	975 Fire Oak, Humble, TX	M	38000	333445555	5
	Joyce	A	English	453453453	31-JUL-62	5631 Rice, Houston, TX	F	25000	333445555	5

R	FIRSTNAME	LASTNAME	SALARY
	John	Smith	30000
	Franklin	Wong	40000
	Ramesh	Narayan	38000
	Joyce	English	25000

## The JOIN Operation

- JOIN denoted by  $\bowtie$  is used to combine related tuples from two relations into single tuples

DEPT\_MGR  $\leftarrow$  DEPARTMENT  $\bowtie_{MGRSSN=SSN}$  EMPLOYEE

RESULT  $\leftarrow \pi$  DNAME, LNAME, FNAME (DEPT\_MGR)

DEPT_MGR	DNAME	DNUMBER	MGRSSN	...	FNAME	MINIT	LNAME	SSN	...
Research	5	333445555	...	Franklin	T	Wong	333445555	...	
Administration	4	987654321	...	Jennifer	S	Wallace	987654321	...	
Headquarters	1	888665555	...	James	E	Borg	888665555	...	

## JOIN

- THETA JOIN
  - comparison operators can be {=, <, >, >=, <=, !=}
- EQUIJOIN
  - the allowed comparison is "="
- NATURAL JOIN
  - EQUIJOIN followed by deleting superfluous attributes
- CROSS JOIN: CARTESIAN PRODUCT
  - no <join condition> is specified

## The NATURAL JOIN operation

(a) PROJ\_DEPT  $\leftarrow$  PROJECT \* DEPT

PROJ_DEPT	PNAME	PNUMBER	PLOCATION	DNUM	DNAME	MGRSSN	MGRSTARTDATE
ProductX	1	Bellaire	5	Research	333445555	22-MAY-78	
ProductY	2	Sugarland	5	Research	333445555	22-MAY-78	
ProductZ	3	Houston	4	Research	333445555	22-MAY-78	
Computerization	10	Stafford	4	Administration	987654321	01-JAN-85	
Reorganization	20	Houston	1	Headquarters	888665555	19-JUN-71	
Newbenefits	30	Stafford	4	Administration	987654321	01-JAN-85	

(b) DEPT\_LOCS  $\leftarrow$  DEPARTMENT \* DEPT\_LOCATIONS

DEPT_LOCS	DNAME	DNUMBER	MGRSSN	MGRSTARTDATE	LOCATION
Headquarters	1	888665555	19-JUN-71	Houston	
Administration	4	987654321	01-JAN-95	Stafford	
Research	5	333445555	22-MAY-78	Bellaire	
Research	5	333445555	22-MAY-78	Sugarland	
Research	5	333445555	22-MAY-78	Houston	

## Aggregate Functions

- Specify mathematical aggregate functions on sets of values from the database
  - e.g. group employee tuples by DNO, and then calculate the average salary of this group
  - SUM, AVERAGE, MAXIMUM, MINIMUM, COUNT
- <grouping attributes>  $\overline{\int}$  <function list> (<relation name>)
  - R(DNO, NUMBER\_OF\_EMPLOYEES, AVERAGE\_SAL)  $\leftarrow$  DNO  $\overline{\int}$  COUNT SSN, AVERAGE SALARY(EMPLOYEE)

## The FUNCTION Operation

(a) R (DNO, NO\_OF\_EMPLOYEES, AVERAGE\_SAL)  $\leftarrow$  DNO  $\overline{\int}$  COUNT SSN, AVERAGE SALARY (EMPLOYEE).

R	DNO	NO_OF_EMPLOYEES	AVERAGE_SAL
5	4	33250	31000
4	3	55000	
1	1	55000	

(b) DNO  $\overline{\int}$  COUNT SSN, AVERAGE SALARY (EMPLOYEE).

DNO	COUNT_SSN	AVERAGE_SALARY
5	4	33250
4	3	31000
1	1	55000

(c)  $\overline{\int}$  COUNT SSN, AVERAGE SALARY (EMPLOYEE).

COUNT_SSN	AVERAGE_SALARY
8	35125

## Two-Level Recursion

(Borg's SSN is 88866555) (Supervised by Borg's subordinates)

SUPERVISION	SSN1	SSN2
	123456789	333445555
	333445555	888665555
	999887777	987654321
	987654321	888665555
	666884444	333445555
	453453453	333445555
	987987987	987654321
	888665555	null

  

(Supervised by Borg)

RESULT1	SSN
	333445555
	987654321

  

(Supervised by Borg's subordinates)

RESULT2	SSN
	123456789
	999887777
	666884444
	453453453
	987987987

  

(Supervised by Borg)

RESULT	SSN
	123456789
	999887777
	666884444
	453453453
	987987987
	333445555
	987654321

(RESULT1)  $\cup$  (RESULT2)

## The LEFT OUTER JOIN operation

RESULT	FNAME	MINIT	LNAME	DNAME
	John	B	Smith	null
	Franklin	T	Wong	Research
	Alicia	J	Zelaya	null
	Jennifer	S	Wallace	Administration
	Ramesh	K	Narayan	null
	Joyce	A	English	null
	Ahmad	V	Jabbar	null
	James	E	Borg	Headquarters

## OUTER UNION

- STUDENT(Name, SSN, Department, Advisor)
- FACULTY(Name, SSN, Department, Rank)
- RESULT(Name, SSN, Department, Advisor, Rank)  
 $\leftarrow$  STUDENT  $\bowtie$  OUTER UNION  $\bowtie$  FACULTY

## Example QUERY 1

Retrieve the name and address of all employees who work for the 'Research' department.

```
RESEARCH_DEPT  $\leftarrow$   $\sigma_{\text{DNAME}='Research'}$  (DEPARTMENT)
RESEARCH_DEPT_EMPS  $\leftarrow$  (RESEARCH_DEPT  $\bowtie_{\text{DNUMBER=DNO}}$  EMPLOYEE)
RESULT  $\leftarrow$   $\pi_{\text{FNAME, LNAME, ADDRESS}}$  (RESEARCH_DEPT_EMPS)
```

This query could be specified in other ways; for example, the order of the JOIN and SELECT operations could be reversed, or the JOIN could be replaced by a NATURAL JOIN.

## Example QUERY 2

- For every project located in 'Stafford', list the project number, the controlling department number, and the department manager's last name, address, and birthdate.

```
STAFFORD_PROJS  $\leftarrow$   $\sigma_{\text{LOCATION}='Stafford'}$  (PROJECT)
CONTR_DEPT  $\leftarrow$  (STAFFORD_PROJS  $\bowtie_{\text{DNUM=DNUMBER}}$  DEPARTMENT)
PROJ_DEPT_MGR  $\leftarrow$  (CONTR_DEPT  $\bowtie_{\text{MGRSSN=SSN}}$  EMPLOYEE)
RESULT  $\leftarrow$   $\pi_{\text{PNUMBER, DNUM, LNAME, ADDRESS, BDATE}}$  (PROJ_DEPT_MGR)
```

## Example QUERY 3

- Find the names of employees who work on *all* the projects controlled by department number 5.

```
DEPT5_PROJS (PNO)  $\leftarrow$   $\pi_{\text{PNUMBER}}$  ( $\sigma_{\text{DNUM}=5}$  (PROJECT))
EMP_PROJ (SSN, PNO)  $\leftarrow$   $\pi_{\text{SSN, PNO}}$  (WORKS_ON)
RESULT_EMP_SSNS  $\leftarrow$  EMP_PROJ  $\div$  DEPT5_PROJS
RESULT  $\leftarrow$   $\pi_{\text{LNAME, FNAME}}$  (RESULT_EMP_SSNS * EMPLOYEE)
```

## Example QUERY 4

- Make a list of project numbers for projects that involve an employee whose last name is 'Smith', either as a worker or as a manager of the department that controls the project.

```
SMITHS(ESSN)  $\leftarrow$   $\pi_{\text{SSN}}$  ( $\sigma_{\text{LNAME}='Smith'}$  (EMPLOYEE))
SMITH_WORKER_PROJS  $\leftarrow$   $\pi_{\text{PNO}}$  (WORKS_ON * SMITH)
MGRS  $\leftarrow$   $\pi_{\text{LNAME, DNUMBER}}$  (EMPLOYEE  $\bowtie_{\text{SSN=MGRSSN}}$  DEPARTMENT)
SMITH_MGRS  $\leftarrow$   $\sigma_{\text{LNAME}='Smith'}$  (MGRS)
SMITH_MANAGED_DEPTS(DNUM)  $\leftarrow$   $\pi_{\text{DNUMBER}}$  (SMITH_MGRS)
SMITH_MGR_PROJS(PNO)  $\leftarrow$   $\pi_{\text{PNUMBER}}$  (SMITH_MANAGED_DEPTS * PROJECT)
RESULT  $\leftarrow$  (SMITH_WORKER_PROJS  $\cup$  SMITH_MGR_PROJS)
```

## Example QUERY 5

- Retrieve the names of employees who have no dependents.

```
ALL_EMPS  $\leftarrow$   $\pi_{\text{SSN}}$  (EMPLOYEE)
EMP_WITH_DEPS (SSN)  $\leftarrow$   $\pi_{\text{ESSN}}$  (DEPENDENT)
EMPS_WITHOUT_DEPS  $\leftarrow$  (ALL_EMPS - EMP_WITH_DEPS)
RESULT  $\leftarrow$   $\pi_{\text{LNAME, FNAME}}$  (EMPS_WITHOUT_DEPS * EMPLOYEE)
```

## Example QUERY 6

- Find the names of managers who have at least one dependent.

```

MGR (SSN) ← πMGRSSN (DEPARTMENT)
EMP_WITH_DEPS (SSN) ← πESSN (DEPENDENT)
MGRS_WITH_DEPS ← (MGRS ∩ EMPS_WITH_DEPS)
RESULT ← πLNAME, FNAME (MGRS_WITH_DEPS * EMPLOYEE)
    
```

## Join conditions for materializing the Relationship types of the COMPANY ER schema

ER Relationship	Participating Relations	Join Condition
WORKS_FOR	EMPLOYEE, DEPARTMENT	EMPLOYEE.DNO = DEPARTMENT.DNUMBER
NAMAGES	EMPLOYEE, DEPARTMENT	EMPLOYEE.SSN = DEPARTMENT.MGRSSN
SUPERVISION	EMPLOYEE(E), EMPLOYEE(S)	EMPLOYEE(E).SUPERSSN = EMPLOYEE(S).SSN
WORKS_ON	EMPLOYEE, WORKS_ON, PROJECT	EMPLOYEE.SSN = WORKS_ON.ESSN AND PROJECT.PNUMBER = WORKS_ON.PNO
CONTROLS	DEPARTMENT, PROJECT	DEPARTMENT.DNUMBER = PROJECT.DNUM
DEPENDENTS_OF	EMPLOYEE, DEPENDENT	EMPLOYEE.SSN = DEPENDENT.ESSN

## Mapping of ER and Relational Models

ER Model	Relational Model
entity type	"entity" relation
1:1 or 1:N relationship type	foreign key (or "relationship" relation)
M:N relationship type	"relationship" relation and two foreign keys
n-ary relationship type	"relationship" relation and n foreign keys
simple attribute	attribute
composite attribute	set of emple component attributes
multivalued attribute	relation and foreign key
value set	domain
key attribute	primary (or secondary) key

## Operations of the Relational Algebra

Operation	Purpose	Notation
SELECT	Selects all tuples that satisfy the selection condition from a relation R.	$\sigma_{\langle \text{selection condition} \rangle} (R)$
PROJECT	Produces a new relation with only some of the attributes of R, and removes duplicate tuples.	$\pi_{\langle \text{attribute list} \rangle} (R)$
THETA JOIN	Produces all combinations of tuples from $R_1$ and $R_2$ that satisfy the join condition.	$R_1 \bowtie_{\langle \text{join condition} \rangle} R_2$
EQUIJOIN	Produces all the combinations of tuples from $R_1$ and $R_2$ that satisfy a join condition with only equality comparisons.	$R_1 \bowtie R_2$ $\text{or } \bowtie_{\langle \text{join condition} \rangle}$
NATURAL JOIN	Same as EQUIJOIN except that the join attributes of $R_2$ are not included in the resulting relation; if the join attributes have the same names, they do not have to be specified at all.	$R_1 \ltimes R_2$ $\text{or } R_1 \ltimes_{\langle \text{join attributes } 1 \rangle} R_2$ $\text{or } R_1 \ltimes_{\langle \text{join attributes } 2 \rangle} R_2$ $\text{or } R_1 \ltimes R_2$

## Operations of the Relational Algebra (Continued)

Operation	Purpose	Notation
UNION	Produces a relation that includes all the tuples in $R_1$ or $R_2$ or both $R_1$ and $R_2$ ; $R_1$ and $R_2$ must be union compatible.	$R_1 \cup R_2$
INTERSECTION	Produces a relation that includes all the tuples in both $R_1$ and $R_2$ ; $R_1$ and $R_2$ must be union compatible.	$R_1 \cap R_2$
DIFFERENCE	Produces a relation that includes all the tuples in both $R_1$ and $R_2$ ; $R_1$ and $R_2$ must be union compatible.	$R_1 - R_2$
CARTESIAN PRODUCT	Produces a relation that has the attributes of $R_1$ and $R_2$ and includes as tuples all possible combinations of tuples from $R_1$ and $R_2$ .	$R_1 \times R_2$
DIVISION	Produces a relation $R(X)$ that includes all tuples $t(X)$ in $R_1(Z)$ that appear in $R_1$ in combination with every tuple from $R_2(Y)$ , where $Z = X \cup Y$ .	$R_1 \div R_2$